

# **Hilbert II**

Presentation of  
Formal Correct  
Mathematical Knowledge

Logical Language

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The source for this document can be found here:

[http://qedeq.org/0\\_03\\_07/doc/project/qedeq\\_logic\\_language.xml](http://qedeq.org/0_03_07/doc/project/qedeq_logic_language.xml)

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# Description

The project **Hilbert II** includes formal correct mathematical knowledge. Here we introduce the underlying formal language for the mathematical formulas. This is done in an informal way. Important theorems (e.g.: universal decomposition, and any proofs) are left out.

All we will do is manipulate symbols. We build lists of symbol strings and use certain simple rules to get new lists. So by starting with a few basic lists we create a whole universe of derived symbol lists. It turns out that these lists could be interpreted as a view to the incredible world of mathematics.



# Chapter 1

## Entities

To describe the logical language we firstly deal with a more basic notation. This notation enables us to formulate the syntax of formulas and terms later on.

### 1.1 Elements, Atoms and Lists

The basic structure we have to deal with is an element. An element is either an atom or a list.

An atom carries textual data, atoms are just strings.

Each list has an operator and can contain elements again. An operator is also nothing more than a simple string. A list has a size: the number of elements it contains. Their elements can be accessed by their position number. An atom has no operator, no size and no subelements in the previous sense.

### 1.2 List Notation

Lists and atoms can be written in the following manner. We write down string atoms quoted with " and the lists as the contents of the operator string followed by ( and a comma separated list of elements and an closing ).

### 1.3 Examples

In this syntax we can write down the following element examples.

```
"I am a string atom"  
EMPTY_LIST()  
THIS_LIST("contains", "three", "atoms")  
OPERATOR("argument 1", "argument 2")  
FUNCTION_A(FUNCTION_B("1", "2"), "3")
```

In the last example we have a list that has the operator FUNCTION\_A and contains two elements. The first element is FUNCTION\_B("1", "2") which is a list too. The second element is the atom "3".



# Chapter 2

## Logical Language

There are different basic things we have to do with. These are predicates, functions, subject variables and logical connectives. In the following all of them are named and described.

### 2.1 Logical Operator Overview

Lists are categorized according to their operators. Before we introduce the formal language in detail the used operators are briefly listed.

<i>logical</i>		
<i>AND</i>	logical conjunction operator	$\wedge$
<i>OR</i>	logical disjunction operator	$\vee$
<i>IMPL</i>	logical implication operator	$\rightarrow$
<i>EQUI</i>	logical biconditional operator	$\leftrightarrow$
<i>NOT</i>	logical negation operator	$\neg$

  

<i>logical quantifiers</i>		
<i>FORLL</i>	universal quantifier	$\forall$
<i>EXISTS</i>	existential quantifier	$\exists$
<i>EXISTSU</i>	unique existential quantifier	$\exists!$

  

<i>variables</i>		
<i>VAR</i>	subject variables	$x, y, z, \dots$
<i>PREDVAR</i>	predicate variables	$A, B, R, \dots$
<i>FUNCVAR</i>	function variables	$f, g, h, \dots$

  

<i>constants</i>		
<i>PREDCON</i>	predicate constants	$=, \in, \subseteq, \dots$
<i>FUNCCON</i>	function constants	$\emptyset, \mathfrak{P}, \dots$
<i>CLASS</i>	class term	$\{x   \phi(x)\}$

### 2.2 Terms and Formulas

Now we define recursively our formal language. We call some elements *subject variables*, *terms* and some other *formulas*. We also define the relations a subject variable *is free in* and *is bound in* a term or a formula. If something is not according to the formal rules errors occur. The error codes are also described.

### 2.2.1 General Error Codes

The atoms and lists that build up a formula or term are subject to restrictions. The following errors occur if an atom has no content or has content with length of 0 or an list has no operator or one of its sub-elements does not exist. These are mainly technical error codes, only the error code 30470 shows an semantical error.

30400	no element	an element doesn't exist - it is null
30410	no atom	an atom doesn't exist - it is null
30420	no list	a list doesn't exist - it is null
30430	no atom content	an atom has no content - it is null
30440	atom content empty	an atom has content with 0 length
30450	no operator	a list has no operator - it is null
30460	operator empty	a list has an operator with 0 length
30470	list expected	list element expected but not found

### 2.2.2 Subject Variable

We call an element *subject variable* iff it has the operator *VAR* and its list size is 1 with an atom as its only argument.

Each subject variable is also called a *term*. Only the subject variable itself is free in itself. No subject variable is bound in a subject variable.

30710	not exactly one argument	list has not exactly one element
30730	atom element expected	the first and only list element must be an atom

### 2.2.3 Function Term

If an element has the operator *FUNVAR* or *FUNCON* and its list size is greater than or equal to 1 with an atom as its first argument and the remaining arguments are all terms then it is called a term too.

Iff a subject variable is free in any sub-element it is also free in the new term. No other subject variables are free. Analogous for bound subject variables.

30720	argument(s) missing	if operator is <i>FUNCON</i> the list must have at least one element
30730	atom element expected	the first list element must be an atom
30740	argument(s) missing	if operator is <i>FUNVAR</i> the list must have more than one element
30770	free bound mixed	found a bound subject variable that is already free in a previous list element
30780	free bound mixed	found a free subject variable that is already bound in a previous list element
30690	undefined constant	the operator is <i>FUNCON</i> and this function constant has not been defined for this argument number

Any other error for term checks may occur due to the fact that all (but the first) sub-elements must be terms too.

### 2.2.4 Predicate Formula

If an element has the operator *PREDVAR* or *PREDCON* and its list size is greater than or equal to 1 with an atom as its first argument and the remaining arguments are all terms and no errors occur then it is called a *formula*.

Iff a subject variable is free in any sub-element it is also free in the new formula. No other subject variables are free. Analogous for bound subject variables.

30720	argument(s) missing	list must have at least one element
30730	atom element expected	the first list element must be an atom
30770	free bound mixed	found a bound subject variable that is already free in a previous list element
30780	free bound mixed	found a free subject variable that is already bound in a previous list element
30590	undefined constant	the operator is <i>PREDCON</i> and this predicate constant has not been defined for this argument number

Any other error for formula checks may occur due to the fact that all (but the first) sub-elements must be terms.

### 2.2.5 Logical Connectives

If an element has the operator *AND*, *OR*, *IMPL* or *EQUI* and its list size is greater than or equal to 2 and the remaining arguments are all formulas and no errors occur then it is called a formula too.

Iff a subject variable is free in any sub-element it is also free in the new formula. No other subject variables are free. Analogous for bound subject variables.

30740	argument(s) missing	list must have more than one element
30760	exactly 2 elements expected	the operator is <i>IMPL</i> and this list size is not equal to 2
30770	free bound mixed	found a bound subject variable that is already free in a previous list element
30780	free bound mixed	found a free subject variable that is already bound in a previous list element

Any other error for formula checks may occur due to the fact that all sub-elements must be formulas.

### 2.2.6 Negation

If an element has the operator *NOT*, its list size is exactly 1 and its only sub-element arguments is a formula then it is called a formula too.

Iff a subject variable is free in the sub-element it is also free in the new formula. No other subject variables are free. Analogous for bound subject variables.

30710	exactly 1 argument expected	list must have exactly than one element
-------	-----------------------------	---

Any other error for formula checks may occur due to the fact that the sub-element must be a formula.

### 2.2.7 Quantifiers

If an element has the operator *FORALL*, *EXISTS* or *EXISTSU* its first sub-element is a subject variable and its second and perhaps its third sub-element is a formula then the element is called a *formula* too.

Iff a subject variable is free in the sub-element it is also free in the new formula. No other subject variables are free. Analogous for bound subject variables.

30760	2 or 3 arguments expected	list must have exactly 2 or 3 elements
30540	subject variable expected	first sub-element must be a subject variable
30550	already bound	subject variable already bound in second or third sub-element
30770	free bound mixed	found a bound subject variable that is already free in a previous list element
30780	free bound mixed	found a free subject variable that is already bound in a previous list element

Any other error for formula checks may occur due to the fact that the sub-element must be a formula.

### 2.2.8 Class Term

An list element with the operator *CLASS*, containing an subject variable and an formula is a term.

Iff a subject variable is free in the formula and is not equal to the first sub-element (which is a subject variable) it is also free in the new term. No other subject variables are free. If a subject variable is bound in the formula it is bound in the new term. Also the first sub-element is bound. No other subject variables are bound.

30760	2 arguments expected	the list must contain exactly two arguments
30540	subject variable expected	the first sub-element must be a subject variable
30550	already bound	the subject variable is already bound in the formula
30680	undefined class operator	the class operator is still unknown

Any other error for formula checks may occur due to the fact that the second sub-element must be a formula.

### 2.2.9 Term

When checking an element for beeing a term the element must have the operator for a *Subject Variable*, *Function Term* or *Class Term*.

30620	unknown term operator	element has no operator that is known as a term operator
-------	-----------------------	--

Any other error for the accordant operator checks may occur.

### 2.2.10 Formula

When checking an element for beeing a formul the element must have the operator for a *Predicate Formula*, *Logical Connective*, *Negation* or *Quantifier*.

30530	unknown logical operator	element has no known logical operator
-------	--------------------------	---------------------------------------

Any other error for the accordant operator checks may occur.

# Chapter 3

## Representations

The representation of elements differ according to the viewpoint. Lets take the following formula for example.

$$y = \{x \mid \phi(x)\} \leftrightarrow \forall z (z \in y \leftrightarrow z \in \{x \mid \phi(x)\})$$

The predicate constant  $\in$  must have been defined in previous sections.

### 3.1 List Notation

In list notation (see 1.2) the above formula looks like the following.

```
EQUI(
  PREDCON(
    "equal",
    VAR("y"),
    CLASS(
      VAR("x"),
      PREDVAR(
        "\phi",
        VAR("x"))
    )
  )
),
FORALL(
  VAR("z"),
  EQUI(
    PREDCON(
      "in",
      VAR("z"),
      VAR("y")
    ),
    PREDCON(
      "in",
      VAR("z"),
      CLASS(
        VAR("x"),
        PREDVAR(
          "\phi",
          VAR("x"))
      )
    )
  )
)
```

)  
)  
)  
)  
)  
)

Due to XSD restrictions for the XML document some error codes listed in Chapter will not occur. Instead the XML will be classified as invalid.

## 3.2 Java

The list notation leads directly to the following Java code.

```
Element el = new ElementListImpl("EQUI", new Element[] {
    new ElementListImpl("PREDCON", new Element[] {
        new AtomImpl("equal"),
        new ElementListImpl("VAR", new Element[] {
            new AtomImpl("y"),
        }),
        new ElementListImpl("CLASS", new Element[] {
            new ElementListImpl("VAR", new Element[] {
                new AtomImpl("x"),
            }),
            new ElementListImpl("PREDVAR", new Element[] {
                new AtomImpl("\u03c6"),
                new ElementListImpl("VAR", new Element[] {
                    new AtomImpl("x"),
                })
            })
        })
    })
}),
new ElementListImpl("FORALL", new Element[] {
    new ElementListImpl("VAR", new Element[] {
        new AtomImpl("z"),
    }),
    new ElementListImpl("EQUI", new Element[] {
        new ElementListImpl("PREDCON", new Element[] {
            new AtomImpl("in"),
            new ElementListImpl("VAR", new Element[] {
                new AtomImpl("z"),
            })
        }),
        new ElementListImpl("VAR", new Element[] {
            new AtomImpl("y"),
        })
    })
}),
new ElementListImpl("PREDCON", new Element[] {
    new AtomImpl("in"),
    new ElementListImpl("VAR", new Element[] {
        new AtomImpl("z"),
    })
}),
new ElementListImpl("CLASS", new Element[] {
    new ElementListImpl("VAR", new Element[] {
        new AtomImpl("x"),
    })
}),
```

```

        new ElementListImpl("PREDVAR", new Element[] {
            new AtomImpl("\phi"),
            new ElementListImpl("VAR", new Element[] {
                new AtomImpl("x"),
            })
        })
    })
}
);
});
```

### 3.3 XML

The XML representation within an QEDEQ module looks a little bit different. Here all first list atoms are represented as the attribute `ref` or `id`. So the above formula may look like the following.

```

<EQUI>
  <PREDCON ref="equal">
    <VAR id="y"/>
    <CLASS>
      <VAR id="x"/>
      <PREDVAR id="\phi">
        <VAR id="x"/>
      </PREDVAR>
    </CLASS>
  </PREDCON>
  <FORALL>
    <VAR id="z"/>
    <EQUI>
      <PREDCON ref="in">
        <VAR id="z"/>
        <VAR id="y"/>
      </PREDCON>
      <PREDCON ref="in">
        <VAR id="z"/>
        <CLASS>
          <VAR id="x"/>
          <PREDVAR id="\phi">
            <VAR id="x"/>
          </PREDVAR>
        </CLASS>
      </PREDCON>
    </EQUI>
  </FORALL>
</EQUI>
```

Due to XSD restrictions for the XML document some error codes listed in Chapter will not occur. Instead the XML will be classified as invalid.

